Advanced Computer Networks

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2009/2010 Semester II

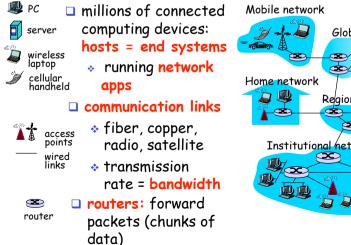
1-1 Review

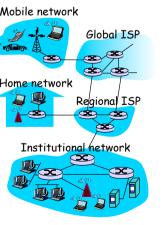
Review on Basic Concepts of Networking

March 11th 2010

Review 1-2

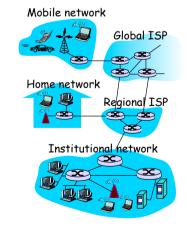
What's the Internet





What's the Internet

- protocols control sending, receiving of msgs ✤ e.g., TCP, IP, HTTP, Skype,
 - Ethernet
- □ Internet: "network of networks"
 - loosely hierarchical
 - public Internet versus private intranet
- Internet standards
 - RFC: Request for comments
 - ✤ IETF: Internet Engineering Task Force



What's a protocol?

network protocols:

 all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

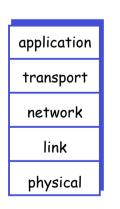
Review 1-5

Internet protocol stack

- application: supporting network applications
 - ✤ FTP, SMTP, HTTP
- transport: process-process data transfer

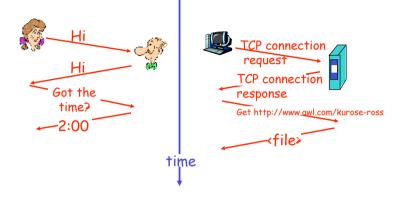
TCP, UDP

- network: routing of datagrams from source to destination
 - $\, \star \,$ IP, routing protocols
- link: data transfer between neighboring network elements
 PPP, Ethernet
- physical: bits "on the wire"



What's a protocol?

a human protocol and a computer network protocol:

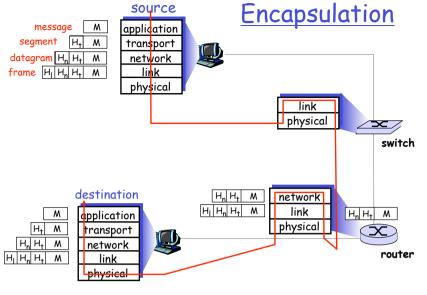


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ISO/OSI reference model

- presentation: allow applications to interpret meaning of data, e.g., encryption, compression, machinespecific conventions
- session: synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
 - these services, if needed, must be implemented in application
 - needed?





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A closer look at network structure:



hosts

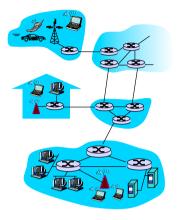
access networks,

physical media: wired, wireless communication links

network core:

 interconnected routers

 network of networks



Review 1-10

The network edge:

end systems (hosts):

- run application programs
- * e.g. Web, email
- at "edge of network"

□ client/server model

- client host requests, receives service from always-on server
- e.g. Web browser/server; emailent/server client/server

□ peer-peer model:

- minimal (or no) use of dedicated servers
- * e.g. Skype, BitTorrent

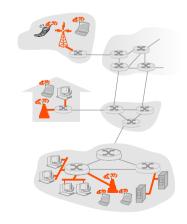


Access networks and physical media

- Q: How to connect end systems to edge router?
- residential access nets
- institutional access networks (school, company)
- mobile access networks

Keep in mind:

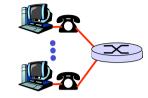
- bandwidth (bits per second) of access network?
- □ shared or dedicated?



Residential access: point to point access

Dialup via modem

- up to 56Kbps direct access to router (often less)
- Can't surf and phone at same time: can't be "always on"



DSL: digital subscriber line

- * deployment: telephone company (typically)
- * up to 1 Mbps upstream
- * up to 28 Mbps downstream
- * dedicated physical line to telephone central office

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Cable Network Architecture: Overview

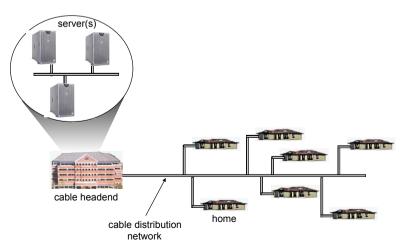


□ HFC: hybrid fiber coax

- asymmetric: up to 30Mbps downstream, 2
 Mbps upstream
- network of cable and fiber attaches homes to ISP router
 - * homes share access to router
- deployment: available via cable TV companies

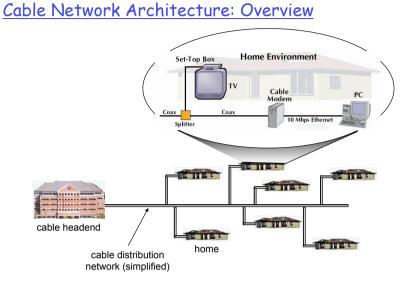
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Cable Network Architecture: Overview



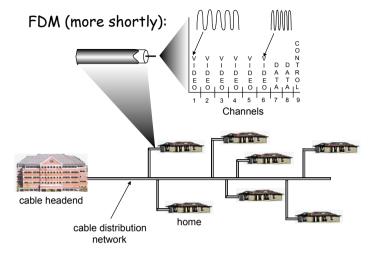
cable distribution network (simplified)

Typically 500 to 5,000 homes



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Cable Network Architecture: Overview



router

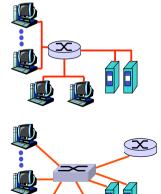
base

station

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Company access: local area networks

- company/univ local area network (LAN) connects end system to edge router
- □ Ethernet:
 - 10 Mbs, 100Mbps, 1Gbps, 10Gbps Ethernet
 - modern configuration: end systems connect into Ethernet switch



Wireless access networks

- □ shared wireless access
 - network connects end system
 - to router
 - via base station aka "access point"

• wireless LANs:

* 802.11b/g (WiFi): 11 or 54 Mbps

wider-area wireless access

- provided by telco operator
- ~1Mbps over cellular system (EVDO, HSDPA)
- next up (?): WiMAX (10's Mbps) over wide area

mobile

hosts

Application architectures

Client-server

Peer-to-peer (P2P)
Hybrid of client-server and P2P

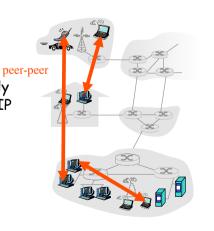
□ specific protocols:

- . ♦ HTTP
- * FTP
- * SMTP, POP, IMAP
- * DNS
- ✤ P2P: BitTorrent, Skype

<u>Pure P2P architecture</u>

- no always-on server
- arbitrary end systems directly communicate
- peers are intermittently connected and change IP addresses

Highly scalable but difficult to manage



<u>Client-server architecture</u>



server:

- always-on host
- * permanent IP address
- server farms for scaling

clients:

- \diamond communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other
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<u>App-layer protocol defines</u>

- Types of messages exchanged,
 - e.g., request, response
- Message syntax:
 - what fields in messages & how fields are delineated
- Message semantics
 - meaning of information in fields
- Rules for when and how processes send & respond to messages

Public-domain protocols:

- □ defined in RFCs
- allows for interoperability
- □ e.g., HTTP, SMTP
- Proprietary protocols:
- 🗆 e.g., Skype

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Transport service requirements of common apps

Applicatio	on	Data loss	Throughput	Time Sensitive
file transf	fer	no loss	elastic	no
e-m	ail	no loss	elastic	no
Web documer	nts	no loss	elastic	no
real-time audio/vide	eo	loss-tolerant	audio: 5kbps-1Mbps video:10kbps-5Mbps	yes, 100's msec
stored audio/vide		loss-tolerant	same as above	yes, few secs
interactive game		loss-tolerant	few kbps up	yes, 100's msec
instant messagii	ng	no loss	elastic	yes and no

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Internet apps: application, transport protocols

Application	Application layer protocol	Underlying transport protocol
e-mail	SMTP [RFC 2821]	ТСР
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (eg Youtube), RTP [RFC 1889]	TCP or UDP
Internet telephony	SIP, RTP, proprietary (e.g., Skype)	typically UDP

Internet transport protocols services

TCP service:

- connection-oriented: setup required between client and server processes
- reliable transport between sending and receiving process
- flow control: sender won't overwhelm receiver
- congestion control: throttle sender when network overloaded
- does not provide: timing, minimum throughput guarantees, security

UDP service:

- unreliable data transfer between sending and receiving process
- does not provide: connection setup, reliability, flow control, congestion control, timing, throughput guarantee, or security
- <u>Q:</u> why bother? Why is there a UDP?

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TCP: Overview

point-to-point: one sender, one receiver

- reliable, in-order byte steam:
- □ pipelined:
 - TCP congestion and flow control set window size

send & receive buffers



🗅 full duplex data:

RFCs: 793, 1122, 1323, 2018, 2581

- bi-directional data flow in same connection
- MSS: maximum segment size

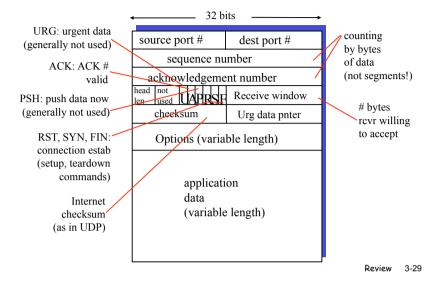
connection-oriented:

 handshaking (exchange of control msgs) init's sender, receiver state before data exchange

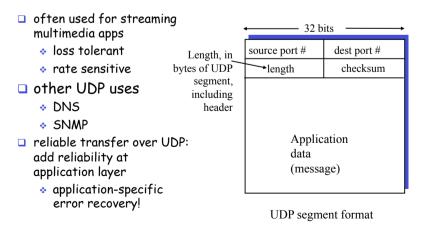
□ flow controlled:

 sender will not overwhelm receiver

TCP segment structure



UDP: more



UDP: User Datagram Protocol [RFC 768]

- "no frills," "bare bones" Internet transport protocol
- "best effort" service, UDP segments may be:
 - lost
 - delivered out of order to app

connectionless:

- no handshaking between UDP sender, receiver
- each UDP segment handled independently of others

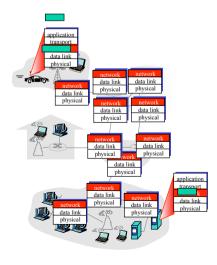
Why is there a UDP?

- no connection establishment (which can add delay)
- simple: no connection state at sender, receiver
- small segment header
- no congestion control: UDP can blast away as fast as desired

Review 3-30

Network layer

- transport segment from sending to receiving host
- on sending side encapsulates segments into datagrams
- on rcving side, delivers segments to transport layer
- network layer protocols in every host, router
- router examines header fields in all IP datagrams passing through it



Two Key Network-Layer Functions

- forwarding: move packets from router's input to appropriate router output
- analogy:
- routing: process of planning trip from source to dest
- routing: determine route taken by packets from source to dest.
 - * routing algorithms
- forwarding: process of getting through single interchange

Review 4-33

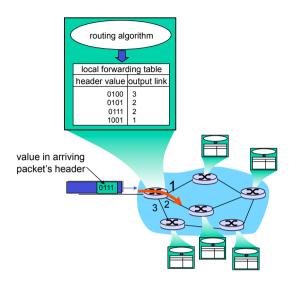
223.1.2

223.1

223.1.3.2

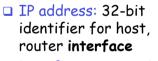
223.1.2.9

Interplay between routing and forwarding



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IP Addressing: introduction



- interface: connection between host/router and physical link
 - router's typically have multiple interfaces
 - host typically has one interface
 - IP addresses associated with each 223.1.1.1 = <u>11011111 00000001 0000001 00000001</u> interface

223.1.1.1

223.1.1.2

223.1.1.4

223.1.1.3 223.1.3.27

223.1.3.1

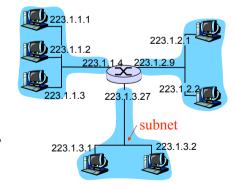
Subnets

□ IP address:

- subnet part (high order bits)
- host part (low order bits)

□ What's a subnet ?

- device interfaces with same subnet part of IP address
- can physically reach each other without intervening router

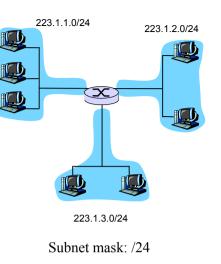


network consisting of 3 subnets

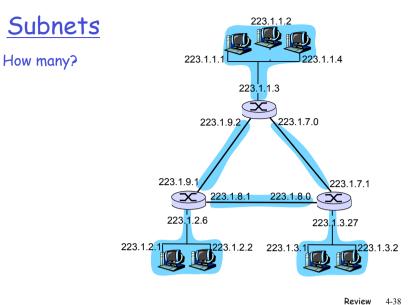
Subnets

Recipe

To determine the subnets, detach each interface from its host or router, creating islands of isolated networks. Each isolated network is called a subnet.







IP addressing: CIDR

CIDR: Classless InterDomain Routing

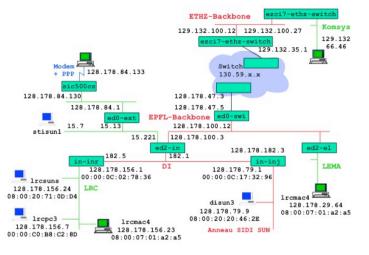
subnet portion of address of arbitrary length
 address format: a.b.c.d/x, where x is # bits in subnet portion of address

subnet <u>host</u> 11001000 00010111 00010000 00000000 200.23.16.0/23

IP addresses: how to get one?

- Q: How does a **host** get IP address?
- □ hard-coded by system admin in a file
 - Windows: control-panel->network->configuration->tcp/ip->properties
 - UNIX: /etc/rc.config
- DHCP: Dynamic Host Configuration Protocol:
 - dynamically get address from as server
 - * "plug-and-play"

Example of Network



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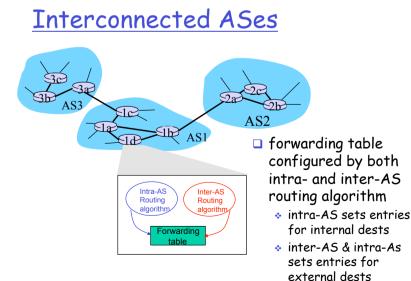
Routing in Internet

- aggregate routers into regions, "autonomous systems" (AS)
- routers in same AS run same routing protocol
 - "intra-AS" routing protocol
 - routers in different AS can run different intra-AS routing protocol

Gateway router

Direct link to router in another AS

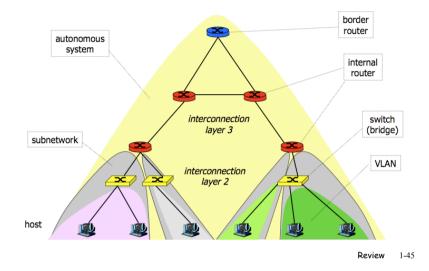
Review 4-42



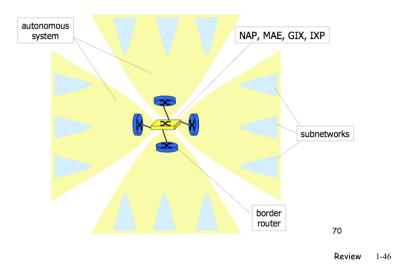
Intra-AS Routing

- also known as Interior Gateway Protocols (IGP)
- most common Intra-AS routing protocols:
 - * RIP: Routing Information Protocol
 - * OSPF: Open Shortest Path First
 - IGRP: Interior Gateway Routing Protocol (Cisco proprietary)

Autonomous Systems



Interconnection of AS

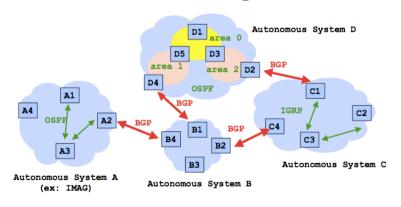


Example of IXP



London Internet Exchange (LINX)

Hierarchical Routing

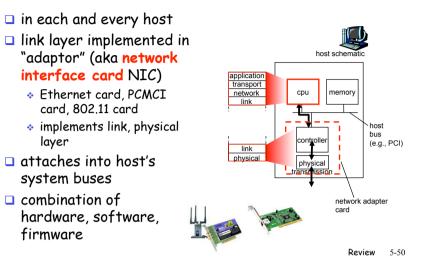


- Hierarchical routing protocols
 - internal (RIP, OSPF, EIGRP)
 - external (BGP)

The Data Link Layer

- Principles behind data link layer services:
 - error detection, correction
 - sharing a broadcast channel: multiple access
 - Iink layer addressing
 - * reliable data transfer, flow control: done!
- link layer technologies

Where is the link layer implemented?



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Multiple Access Links and Protocols

Two types of "links":

point-to-point

- PPP for dial-up access
- * point-to-point link between Ethernet switch and host
- broadcast (shared wire or medium)
 - old-fashioned Ethernet
 - upstream HFC
 - 802.11 wireless LAN







shared wire (e.g., shared RF cabled Ethernet) (e.g., 802.11 WiFi)

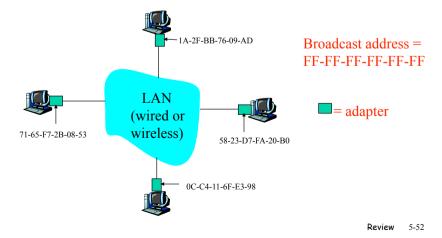


humans at a cocktail party (shared air, acoustical)

LAN Addresses and ARP

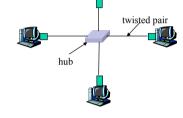
Each adapter on LAN has unique LAN address

layer



Hubs

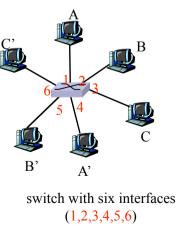
- ... physical-layer ("dumb") repeaters:
 - bits coming in one link go out all other links at same rate
 - \ast all nodes connected to hub can collide with one another
 - no frame buffering
 - no CSMA/CD at hub: host NICs detect collisions



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Review 5-53
```

<u>Switch: allows **multiple** simultaneous</u> <u>transmissions</u>

- hosts have dedicated, direct connection to switch
 awitches buffer peakets
- switches buffer packets
- Ethernet protocol used on each incoming link, but no collisions; full duplex
 - each link is its own collision domain
- switching: A-to-A' and Bto-B' simultaneously, without collisions
 - * not possible with dumb hub



<u>Switch</u>

- link-layer device: smarter than hubs, take active role
 - store, forward Ethernet frames
 - examine incoming frame's MAC address, selectively forward frame to one-or-more outgoing links when frame is to be forwarded on segment, uses CSMA/CD to access segment

transparent

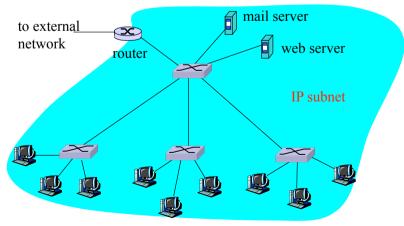
hosts are unaware of presence of switches

plug-and-play, self-learning

* switches do not need to be configured

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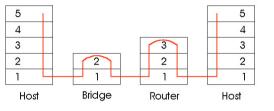
Institutional network



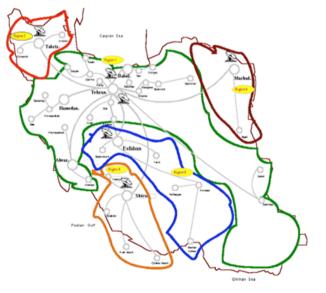
Switches vs. Routers

both store-and-forward devices

- routers: network layer devices (examine network layer headers)
- * switches are link layer devices
- routers maintain routing tables, implement routing algorithms
- switches maintain switch tables, implement filtering, learning algorithms

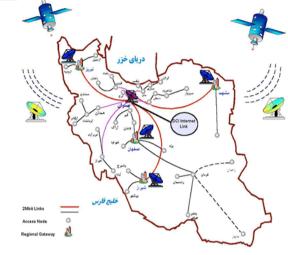


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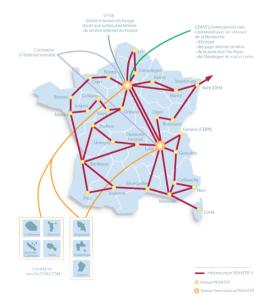
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Scientific Infra. Of Iran

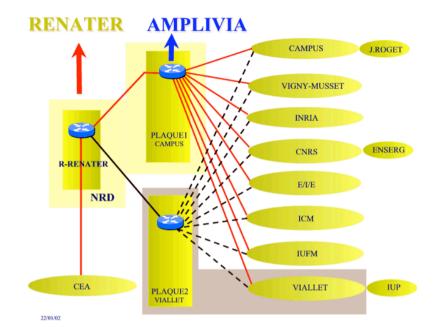


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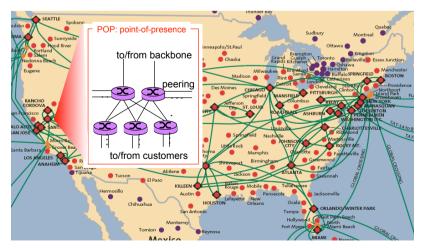
Infrastructure du réseau







<u>Tier-1 ISP: e.g., Sprint</u>



<u>Contents</u>

- Introduction
 - TCP/IP model
- Interconnection Layer 2
 - * VLANs and bridges, spanning tree protocol
- □ Interconnection Layer 3
 - IPv6
 - Routing (RIP, OSPF, BGP)
- Congestion control
- Quality of service

<u>Contents</u>

MPLS, multicast
Mobility
Network management
...

Course Grade

Assignments

- Including simulations, hands-on assignments
- Research Project
 - Choose papers among highly-cited recent published papers in wireless domain.
- Mid-term exam
- Final exam

Review 1-65

Reference Materials

Book Chapters

Based on subject

Research Papers

Standards

Backup Slides

Residential access: cable modems

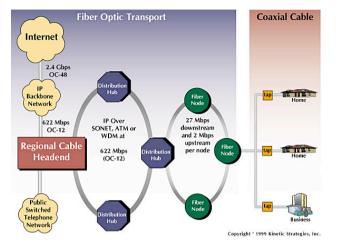


Diagram: http://www.cabledatacomnews.com/cmic/diagram.html